
Online and Approximation Algorithms

Due April 22, 2016 at 08:15, in class

Exercise 1 (Ski Rental - 10 points)

The ski rental problem is defined as follows: Assume that renting a pair of skis costs 1 per day while buying a pair of skis costs b . Every day we have to decide, in an online fashion, whether we want to continue renting skis for another day or buy a pair of skis. At some unknown time D , we will break our leg and have to quit skiing. Our goal is to minimize the cost of skiing.

- What is the optimal offline cost?
- Develop a $(2 - \frac{1}{b})$ -competitive online algorithm ALG for the ski rental problem and prove its competitiveness.

Exercise 2 (List Scheduling on Unrelated Machines - 10 points)

In class, an online greedy algorithm for the problem of scheduling n jobs on m identical machines was presented. In this exercise, we consider the same problem in a more complex environment with m *unrelated* machines M_1, M_2, \dots, M_m .

In this setting, the processing time of a job J_j , $1 \leq j \leq n$, depends on the machine on which it is executed. Specifically, if job J_j is executed by machine M_i , then its processing time is $p_{i,j}$. The greedy algorithm presented in class for identical machines can be easily extended for unrelated machines as follows: Schedule each job J_j on a machine M_i that results in the schedule with minimum makespan. Prove that the competitive ratio of GREEDY is not smaller than $\frac{m}{1 + \varepsilon}$, where ε is an arbitrary small constant.

Hint: There exists an example where each job has finite processing time only on two machines and infinite processing time on all the other machines.

Exercise 3 (First-in First-out - 10 points)

Recall that FIFO is the online paging algorithm that evicts the page that has been in fast memory for the longest time. Prove that FIFO is k -competitive, where k is the number of pages that fit in fast memory.

Exercise 4 (Demand Paging - 10 points)

Paging algorithms that do not evict pages unless there is a page fault are called *demand paging*. Prove that any paging algorithm can be modified to be demand paging without increasing the overall number of memory replacements on any request sequence.